

# bio $\kappa$ : a simple calculus for proteins and cells

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## Some features

- Several agents may interact at the same time by means of several domains (sites)
  - parallelism
  - competition
  - nondeterminism
- The overall behaviour is **deterministic**.
- Interactions may involve simple agents – **proteins** – or complex ones – **cells** – and may cause small local changes or huge structural changes.

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## Two different directions

Two different approaches :

- Based on  $\pi$ -calculus (Regev-Shapiro, Danos-Laneve) :  $\kappa$ -calculus
- Based on Ambients (Cardelli) : Brane Calculi

For modelling different biological systems :

- Signal transduction pathways, gene regulatory networks, ...
- Molecular transport, virus infections, ...

# À la $\pi$ -calcul

We intend to pursue an **algebraic approach** à la pi-calculus

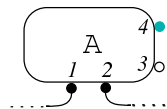
- few (biological) constructs
- a “faithful” rendering of biological interactions – not *via* an encoding
- a **compositional** semantics based on the notion of **interaction**

But simplicity has a cost ! we are loosing :

- expressiveness
- stochastic behaviours

# A calculus for proteins ...

## Proteins



- *visible site*
- *hidden site*
- *bound site*

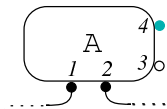
is described in bio $\kappa$  by  $A(1^x + 2^z + 3^v + 4^h)$  – actually we write  $A(1^x + 2^z + 3 + \bar{4})$

syntactically, a protein is  $A(\sigma)$  :

- A belongs to a countable set of *protein names*
- for every A,  $s(a)$  gives an integer – the number of *sites*
- there is a set  $\mathcal{E}$  of *edge names* that are ranged over by  $x, y, z$ , etc.
  - $v, h \notin \mathcal{E}$
- $\sigma$  is a total function from  $1..s(a)$  to  $\{v, h\} \cup \mathcal{E}$  such that  $\sigma$  is injective on  $\mathcal{E}$

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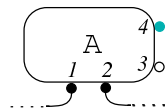
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## ... and cells

cells  $m(M)[S]$

- $m$  belongs to a countable set of *membrane types*
- $M$  is the *membrane*
- $S$  is a biological solution (that may contain cells)
- **well-formedness constraints** :
  - (edge-condition) every solution is such that edge names occur at most twice ;
  - (membrane-condition) every membrane is a multiset of proteins, that is cells do not occur in membranes – *we are abstracting out the bilipidic layer*
  - (nucleus-condition) the dangling edges of nuclei of cells are connected to the corresponding membrane

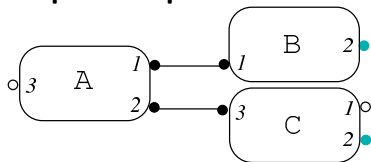
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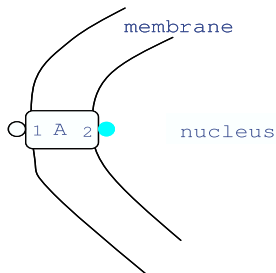
Examples en bio $\kappa$ 

## Complexe of proteins :



$$A(1^x + 2^y + 3), B(1^x + \bar{2}), C(1 + \bar{2} + 3^y)$$

## Cell with a transmembrane protein :



$$m(A(1 + \bar{2})) [S]$$

**remark** : we do not specify whether a site of a protein is outside or inside a membrane

# bio $\kappa$ : The syntax

*Solutions S :*

$S ::=$	<b>solution</b>
$\mathbf{0}$	(empty solution)
$A(\sigma)$	(protein)
$m(S)[S]$	(cell)
$S, S$	(group)

## Some notations

- $v$ - $h$ -maps, ranged over by  $\phi, \psi, \dots$ , are partial maps from naturals to  $\{v, h\}$

$$- \bar{\phi} \text{ is the } v\text{-}h\text{-map : } \bar{\phi}(i) = \begin{cases} h & \text{if } \phi(i) = v \\ v & \text{if } \phi(i) = h \\ \text{undefined} & \text{otherwise} \end{cases}$$

- $\alpha, \beta$ , etc. range over  $(A, a, \phi)$ , such that  $\{a\} \uplus \text{dom}(\phi) \subseteq 1..s(A)$
- complexations**  $\mathcal{C}$  and **decomplexations**  $\mathcal{D}$  are functions from rule names  $r$  to tuples  $(\alpha, \beta)$  with disjoint domain.

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# bioκ : The labelled transition system

The **transition relation**  $\xrightarrow{\mu}$  is the least one satisfying the reductions :

- protein-protein reductions**

$$\frac{(A, a, \phi) \in \mathcal{C}(\mathbf{r}) \quad x \notin \text{en}(\sigma)}{A(a + \phi + \sigma) \xrightarrow{A_r^x} A(a^x + \bar{\phi} + \sigma)} \quad \frac{(A, a, \phi) \in \mathcal{D}(\mathbf{r})}{A(a^x + \phi + \sigma) \xrightarrow{A_r^x} A(a + \bar{\phi} + \sigma)}$$

$$\frac{S \xrightarrow{A_r^x} S' \quad T \xrightarrow{B_r^x} T' \quad A \neq B}{S, T \xrightarrow{\tau} S', T'} \quad \frac{S \xrightarrow{\mu} S' \quad \text{diff}(S, S') \cap \text{en}(T) = \emptyset}{S, T \xrightarrow{\mu} S', T}$$

plus the symmetric rule for groups

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$$\frac{M \xrightarrow{A_r^x} M' \quad S \xrightarrow{B_r^x} S' \quad A \neq B}{m(M)[S] \xrightarrow{\tau} m(M')[S']}$$

- cellular reductions

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remarks 1. every edge name created in a complexation is fresh

2. the interaction between a protein outside a cell and the membrane of the cell is modelled by the last rule and the reaction rule

3. the nucleus of a cell cannot interact with agents external to the cell

4. the reductions do not change the cellular structure – core-bioκ

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- 1 a dimeric form of the epidermal growth factor EGF binds two receptors EGFR located on some plasmic membrane
- 2 the receptors EGFR cross-phosphorylate each other through their tyrosine kinase sites
- 3 then the EGFR activate another binding site that binds an adapter protein SHC and activate it
- 4 ... the signal goes further till reaching the nucleus

$$((\text{EGF}, 1, \bar{2}), (\text{EGF}, 1, \bar{2})) \in \mathcal{C} \quad (1)$$

$$((\text{EGF}, 2, \emptyset), (\text{EGFR}, 1, \bar{4})) \in \mathcal{C} \quad (2)$$

$$((\text{EGFR}, 2, \bar{3} + 4), (\text{EGFR}, 2, \bar{3} + 4)) \in \mathcal{C} \quad (3)$$

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# bio $\kappa$ : The RTK-MAPK pathway

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# To compare the systems

Some notations :

- $S \xRightarrow{\tau} S'$  means  $S \xrightarrow{\tau}^* S'$
- $S \xRightarrow{\mu} S'$ , with  $\mu \neq \tau$ , means  $S \xrightarrow{\tau}^* \xrightarrow{\mu} \xrightarrow{\tau}^* S'$

A (*weak*) *bisimulation* is a symmetric binary relation  $\mathcal{R}$  between solutions such that  $S \mathcal{R} T$  implies :

- 1 if  $S \xrightarrow{\tau} S'$  then  $T \xRightarrow{\tau} T'$  and  $S' \mathcal{R} T'$
- 2 if  $S \xrightarrow{A_r^x} S'$  then  $T \xRightarrow{A_r^x} T'$  and  $S' \mathcal{R} T'$ .

$S \approx T$  if  $S \mathcal{R} T$  for some bisimulation  $\mathcal{R}$ .

# Basic properties

- The transition system preserves the well-formedness constraints

- “, ” is an abelian monoidal operator with identity  $\mathbf{0}$  :

$$S, T \approx T, S \quad (S, T), R \approx S, (T, R) \quad S, \mathbf{0} \approx S$$

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## Merging membranes

- core-bio $\kappa$  preserves the cellular structure of the solution
- it is not possible to describe phenomena such as *fusions* between endosomes :

$$esm(\langle M \rangle [S]) , esm(\langle N \rangle [T]) \longrightarrow esm(\langle M , N \rangle [S , T])$$

**question** : how to define the semantics of an endosome, regardless of endosomes in the context ?

- **answer** : by means of higher-order mechanisms

# The core bio $\kappa$ with mreagents

The syntax of bio $\kappa$  :

$S ::=$	<b>solution</b>
$\mathbf{0}$	(empty solution)
$A(\sigma)$	(protein)
$m \langle M \rangle [S]$	(cell)
$S, S$	(group)
$\langle M ; S \rangle \cdot S$	(mreagents)

**constraint** : in  $\langle M ; S \rangle \cdot T$

- $S$  and  $T$  do not contain mreagents
- $M$  is a multiset of proteins
- $\text{de}(S) \subseteq \text{de}(M)$

# Fusions and activations

- The **fusion function**  $\mathcal{F}$  from rule name to pairs  $(m \otimes m', n)$
- The **activation function**  $\mathcal{A}$  from pairs  $(A_{\mathbf{r}}, m)$  to membrane type.
- We assume that  $\mathcal{C}, \mathcal{D}, \mathcal{F}, \mathcal{A}$  have a disjoint domaine.

$$\frac{m \in \mathcal{F}(\mathbf{r})}{m(M)[S] \xrightarrow{m_{\mathbf{r}}} \langle M; S \rangle \cdot \mathbf{0}}$$

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still to be done :

– properties that reduce the quantifications :

- instead of 4 we may use :

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# Conclusions

- Contribution :
  - one framework for two very different types of biological systems
  - sometimes with a finer description of the phenomena
  - reuse tools already developed for the pi-calculus
- Rules creating cells have been hidden
- The missing part :
  - comparisons with other models (Ambients)
  - more studies on the good notion of bisimulation
  - adding other significant primitives